

Lecture II: Computing Kazhdan-Lusztig Polynomials

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1. Hecke Algebras

- The Hecke Algebra \mathcal{H} associated with a Coxeter system (W, S) is a free $R = \mathbb{Z}[q^{1/2}, q^{-1/2}]$ module with basis $\{T_w \mid w \in W\}$ and the multiplication defined by

$$T_w T_s = \begin{cases} T_{ws}, & \text{if } \ell(ws) > \ell(w), \\ qT_{ws} + (q-1)T_w, & \text{if } \ell(ws) < \ell(w). \end{cases}$$

- As an R -algebra, \mathcal{H} is generated by $\{T_s \mid s \in S\}$ subject to relations:
 - 1 $T_s = T_s T_1 = T_1 T_s$ and $T_s^2 = qT_1 + (q-1)T_s, \forall s \in S$; and
 - 2 $T_{s_1} T_{s_2} \cdots$ (m terms) $= T_{s_2} T_{s_1} \cdots$ (m terms), where $s_1, s_2 \in S$ and $m = o(s_1 s_2)$.
- Note that $T_s, s \in S$, is invertible and

$$T_s^{-1} = q^{-1} T_s - (1 - q^{-1}) T_1.$$

Thus T_w is invertible for all $w \in W$.

1. Hecke Algebras

- Define an involution $\bar{}$ on \mathcal{H} by requiring

$$\overline{q^{1/2}} = q^{-1/2}, \quad \overline{T_w} = (T_{w^{-1}})^{-1},$$

and $\overline{hh'} = \bar{h} \cdot \bar{h'}$, $\forall h, h' \in \mathcal{H}$. Let \mathcal{H}^0 be the set of fixed points of this involution, and let $R^0 = \mathcal{H}^0 \cap R$. Then $R^0 = \mathbb{Z}[f]$ where $f = q^{1/2} + q^{-1/2}$. Note that \mathcal{H}^0 is an R^0 -algebra, and as an R^0 -module it is free.

- Theorem 1.1.** (Kazhdan-Lusztig) For $y \in W$, there is a unique element $C_y \in \mathcal{H}^0$ such that

$$C_y = \sum_{x \leq y} \varepsilon(x)\varepsilon(y)q^{(\ell(y)-2\ell(x))/2}\overline{P}_{x,y}T_x,$$

where $P_{x,y} \in R$ is a polynomial in q of degree $\leq (\ell(y) - \ell(x) - 1)/2$ for $x < y$, and $P_{y,y} = 1$.

2. Kazhdan-Lusztig Polynomials

- The polynomials $P_{x,y}$ are the so-called *Kazhdan-Lusztig polynomials*.
- Theorem 1.1 can be restated as

Theorem 2.1. For $y \in W$, there is a unique element $C'_y \in \mathcal{H}^0$ such that

$$C'_y = q^{-\ell(y)/2} \sum_{x \leq y} P_{x,y} T_x,$$

where $P_{x,y} \in R$ is a polynomial in q of degree $\leq (\ell(y) - \ell(x) - 1)/2$ for $x < y$, and $P_{y,y} = 1$.

- The elements C'_y and C_y are related by $C'_y = (-1)^{\ell(y)} \sigma(C_y)$, where σ is the involution of \mathcal{H} defined by

$$\sigma\left(\sum a_w T_w\right) = \sum (-1)^{\ell(w)} \bar{a}_w q^{-\ell(w)} T_w.$$

2. Kazhdan-Lusztig Polynomials

- **Conjecture 2.2.** (Kazhdan-Lusztig, not “the Kazhdan-Lusztig conjecture”) The coefficients of the polynomials $P_{x,y}$ are non-negative integers.
- The conjecture is true for the Weyl groups of Kac-Moody Lie algebras.
- The values of these polynomials at 1 provide character formulas for the infinite dimensional simple highest weight modules of semisimple Lie algebras over \mathbb{C} (the Kazhdan-Lusztig conjecture, true for symmetrizable Kac-Moody algebras):

$$chL_w = \sum_{y \leq w} \varepsilon(yw) P_{y,w}(1) chM_y.$$

- For real Lie groups, one has the Lusztig-Vogan polynomials.

2. Kazhdan-Lusztig Polynomials

- These polynomials have a rich combinatorial theory, and there have been a lot of research activities devoted to study them.
- **Question:** How to compute these polynomials?
- Kazhdan and Lusztig provided a recursive method in their original paper for the computation of these polynomials via another set of polynomials $R_{x,y}$ defined by

$$\overline{T}_y = (T_{y^{-1}})^{-1} = \sum_x \overline{R}_{x,y} q^{-\ell(x)} T_x.$$

- **Proposition 2.3.** We have
 - $\overline{R}_{x,y} = \varepsilon(xy) q^{\ell(x)-\ell(y)} R_{x,y}.$
 - $\overline{T}_y = \sum_{x \leq y} \varepsilon(xy) q^{-\ell(y)} R_{x,y} T_x.$
 - $\sum_{x \leq z \leq y} \varepsilon(xz) R_{x,z} R_{z,y} = \delta_{x,y}.$

2. Kazhdan-Lusztig Polynomials

- The polynomials $R_{x,y}$ can be computed recursively using the formulas:
 - $R_{x,y} = 1$ if $x = y$ and 0 unless $x \leq y$,
 - $R_{x,y} = R_{sx, sy}$ if $sx < x$ and $sy < y$,
 - $R_{x,y} = R_{xs, ys}$ if $x < xs$ and $y < ys$,
 - $R_{x,y} = (q - 1)R_{sx,y} + qR_{sx, sy}$ if $sx > x$ and $sy < y$.
- The polynomials $P_{x,w}$ can then be constructed using

$$q^{(\ell(w)-\ell(x))/2} \bar{P}_{x,w} - q^{(\ell(x)-\ell(w))/2} P_{x,w} = q^{(\ell(x)-\ell(w))/2} \sum_{x < y \leq w} R_{x,y} P_{y,w}.$$

2. Kazhdan-Lusztig Polynomials

- **Exercise.** Compute $P_{x,y}$ for $y = s_1 s_2 s_1 \in S_3$ using these recursive formulas.
- Next, we describe a closed formula due to Deodhar. The theorems, propositions, etc, in the rest of this lecture are due to Deodhar. Let us first consider an example.
- **Example.** Compute $P_{x,y}$ for $y = s_1 s_2 s_3 s_1 \in S_4$. Recall that $P_{x,y}$ is a polynomial in q of degree $\leq (\ell(y) - \ell(x) - 1)/2$ for $x < y$.

3. Combinatorics of Subexpressions

- Let $y \in W$. Fix a reduced expression $y = s_1 s_2 \cdots s_k$, and let $s = (s_1, s_2, \dots, s_k)$. We give an equivalent definition of the subexpressions of y .
- **Definition 3.1.** We call $\sigma = (\sigma_0, \sigma_1, \dots, \sigma_k)$ a subexpression of s if
 - (1) $\sigma_0 = 1$, and
 - (2) $\sigma_{j-1}^{-1} \sigma_j \in \{1, s_j\}, \forall 1 \leq j \leq k$.
- The set of all subexpressions is denoted by \mathcal{S} .
- Define maps ℓ , π , and $\tilde{\cdot}$ on \mathcal{S} by
 - 1 $\ell(\sigma) = |\{j \mid \sigma_{j-1} \neq \sigma_j\}|$ ($= t$ if σ is identified with $1 \leq i_1 < i_2 < \cdots < i_t \leq k$).
 - 2 $\pi(\sigma) = \sigma_k$.
 - 3 $\tilde{\sigma} = \{(\sigma_0, \dots, \sigma_{k-1}, \sigma_k s_k)\}$.

3. Combinatorics of Subexpressions

- We have

1 $\text{Im}(\pi) = \{x \in W \mid x \leq y\} := W(y).$

2 $\ell(\sigma) - \ell(\pi(\sigma)) \in 2\mathbb{Z}_+.$

3 $\ell(\tilde{\sigma}) = \ell(\sigma) \pm 1.$

- **Definition 3.2.** The *defect* of σ is defined to be

$$d(\sigma) = |\{j \mid \sigma_{j-1}s_j < \sigma_{j-1}\}|.$$

- **Example 3.3.** Consider S_3 . Let $y = s_1s_2s_1$. Then $s = (s_1, s_2, s_1)$ and (since $\sigma_0 = 1$, we omit it in our listing)

$$\mathcal{S} = \{(1, 1, 1), (1, 1, s_1), (1, s_2, 1), (s_1, 1, 1), \\ (1, s_2, s_1), (s_1, 1, s_1), (s_1, s_2, 1), (s_1, s_2, s_1)\}.$$

For $\sigma = (s_1, 1, s_1)$, we have $\ell(\sigma) = 2$, $\pi(\sigma) = 1$, $\tilde{\sigma} = (s_1, 1, 1)$, and $d(\sigma) = 1$.

3. Combinatorics of Subexpressions

■ **Proposition 3.3.** We have

1 $d(\sigma) = d(\tilde{\sigma})$.

2 $\pi : \mathcal{D}_0 = \{\sigma \mid d(\sigma) = 0\} \longrightarrow W(y)$ is bijective.

3 The subset \mathcal{D}_0 is stable under the map \sim .

4 $d(\sigma) \geq (\ell(\sigma) - \ell(\pi(\sigma)))/2$.

5 $\ell(y) - \ell(\pi(\sigma)) \geq d(\sigma)$ and equality holds if and only if $\pi(\sigma) = y$.

■ For $\mathcal{L} \subset \mathcal{S}$, define $h(\mathcal{L}) \in \mathcal{H}$ by

$$h(\mathcal{L}) = q^{-\ell(y)/2} \sum_{\sigma \in \mathcal{L}} q^{d(\sigma)} T_{\pi(\sigma)}.$$

Set

$$P_x(\mathcal{L}) = \sum_{\substack{\sigma \in \mathcal{L} \\ \pi(\sigma) = x}} q^{d(\sigma)}.$$

Then

$$h(\mathcal{L}) = q^{-\ell(y)/2} \sum_{x \leq y} P_x(\mathcal{L}) T_x.$$

4. Admissible Subsets

- **Definition 4.1.** A subset $\mathcal{E} \subset \mathcal{S}$ is called *admissible* if it satisfies the following three properties:

(1) \mathcal{E} contains the special subexpression

$$(1, s_1, s_1 s_2, \dots, s_1 s_2 \cdots s_k) = \pi^{-1}(y),$$

(2) $\mathcal{E} = \tilde{\mathcal{E}}$, and

(3) $h(\mathcal{E}) \in \mathcal{H}^0$.

- Note that \mathcal{S} is admissible. Conditions (1) and (2) clearly hold. (3) can be seen from Proposition 5.1 below.

- **Lemma 4.2.** Let $\mathcal{L} \subset \mathcal{S}$. Then $h(\mathcal{L}) \in \mathcal{H}^0$ if and only if

$$\sum_{x \leq z \leq y} R_{x,z} P_z(\mathcal{L}) = q^{\ell(y) - \ell(x)} \overline{P_x(\mathcal{L})} \quad \forall x \in W(y).$$

- **Lemma 4.3.** If \mathcal{E} is admissible then $\mathcal{D}_0 \subset \mathcal{E}$.

4. Admissible Subsets

- Let R_+^0 be the set of polynomial in $f = q^{1/2} + q^{-1/2}$ with non-negative coefficients. For a pair $\mathcal{E}_1 \supset \mathcal{E}_2$ of admissible sets, we say that \mathcal{E}_1 dominates \mathcal{E}_2 if there is an $x \in W(y)$ and $g \in R_+^0$ such that
 - (1) $h(\mathcal{E}_1 \setminus \mathcal{E}_2) = g \cdot C'_x$, and
 - (2) $\deg(P_z(\mathcal{E}_2)) \leq (\ell(y) - \ell(z) - 1)/2$ for all $x \leq z < y$.
- Let \leq be the partial order generated by this dominance relation. Note that \mathcal{S} is a maximal element for this order.

5. Another Basis of \mathcal{H}^0

- Theorem 2.1 implies that $C'_x, x \in W$ form a basis of the R -module \mathcal{H} and hence a basis of the R^0 -module \mathcal{H}^0 .
- Also from Theorem 2.1, one has $C'_s = q^{-1/2}(T_s + T_1)$ for $s \in S$.
- For $\forall y \in W$, fix a reduced expression $y = s_1 \cdots s_k$. Let $s = (s_1, \dots, s_k)$, and let

$$D'_s = C'_{s_1} C'_{s_2} \cdots C'_{s_k}.$$

- **Proposition 5.1.** (1) We have

$$D'_s = q^{-\ell(y)/2} \sum_{\sigma \in \mathcal{S}} q^{d(\sigma)} T_{\pi(\sigma)}.$$

- (2) The set $\{D'_s \mid y \in W\}$ with s as described above, is a R^0 basis of \mathcal{H}^0 .

5. Another Basis of \mathcal{H}^0

- Let $\mu(x, z)$ be the coefficient of $q^{(\ell(z)-\ell(x)-1)/2}$ in $P_{x,z}$.
- **Proposition 5.2.** Let $s^\# = (s_1, \dots, s_{k-1})$, denote the corresponding terms by using the notation $\#$, and let

$$D'_s = C'_y + \sum_{x < y} L_x C'_x \quad \text{and} \quad D'_{s^\#} = C'_{y_{s_k}} + \sum_{w < y_{s_k}} L_w^\# C'_w,$$

where L_x and $L_w^\# \in R^0$. Then

$$L_x = \begin{cases} 0, & xs_k > x, \\ fL_x^\# + L_{xs_k}^\# + \sum_{\substack{zs_k > z \\ z \leq ys_k}} \mu(x, z)L_z^\#, & xs_k < x. \end{cases}$$

Furthermore, if $\mu(x, z) \geq 0$ for all $x \leq z (\in W(y))$, then L_x is a polynomial in f with non-negative coefficients for all $x \in W(y)$.

6. Minimal Admissible Sets

- **Proposition 6.1.** Let \mathcal{E} be an admissible set. Then there exists a unique subset $A(\mathcal{E})$ of elements strictly less than y such that

$$h(\mathcal{E}) = C'_y + \sum_{x \in A(\mathcal{E})} L_x C'_x.$$

Or equivalently,

$$P_u(\mathcal{E}) = P_{u,y} + \sum_{x \in A(\mathcal{E})} q^{(\ell(y) - \ell(x))/2} L_x P_{u,x}, \quad \forall u \leq y.$$

6. Minimal Admissible Sets

■ **Corollary 6.2.** Let \mathcal{E} be an admissible set.

(a) The following are equivalent.

(1) $A(\mathcal{E}) = \emptyset$.

(2) $h(\mathcal{E}) = C'_y$.

(3) $\deg P_u(\mathcal{E}) \leq (\ell(y) - \ell(u) - 1)/2$ for all $u < y$.

(b) If the conditions of part (a) hold, then \mathcal{E} is minimal for the order \leq .

■ **Definition 6.3.** An element $y \in W$ is called *good* if for any minimal admissible set \mathcal{E}_{min} , $h(\mathcal{E}_{min}) = C'_y$.

6. Minimal Admissible Sets

- **Proposition 6.4.** Let y be good, let \mathcal{E} be an admissible set, and let x_0 be an element such that

$$\deg P_z(\mathcal{E}) \leq \frac{\ell(y) - \ell(z) - 1}{2}, \quad \forall x_0 < z < y,$$

and

$$\deg P_{x_0}(\mathcal{E}) > \frac{\ell(y) - \ell(x_0) - 1}{2}.$$

Then there is an admissible set \mathcal{E}' and $g \in R_+^0$ such that \mathcal{E}' is dominated by \mathcal{E} and $h(\mathcal{E} \setminus \mathcal{E}') = gC'_{x_0}$.

6. Minimal Admissible Sets

- **Theorem 6.5.** If the coefficients of the Kazhdan-Lusztig polynomials $P_{x,z}$ are non-negative for all $x \leq z \leq y$, then y is good and hence

$$P_{x,y} = \sum_{\substack{\sigma \in \mathcal{E}_{min} \\ \pi(\sigma)=x}} q^{d(\sigma)} \quad (= P(\mathcal{E}_{min})),$$

where \mathcal{E}_{min} is any minimal admissible set.

- **Corollary 6.6.** If (W, S) is the Weyl group of a Kac-Moody Lie algebra, then every element of W is good and hence Theorem 6.5 provides a closed formula as well as an algorithm for all the Kazhdan-Lusztig polynomials.

7. Algorithm

- **Algorithm 7.1.** Let y be good and fix a reduced expression

$$y = s_1 \cdots s_k.$$

(1) Start with \mathcal{S} . If

$$\deg P_z(\mathcal{S}) \leq \frac{\ell(y) - \ell(z) - 1}{2}, \quad \forall z < y,$$

then $h(\mathcal{S}) = C'_y$, and hence $P_z(\mathcal{S}) = P_{z,y}$ for all $z \leq y$.

(2) If this is not the case, then \mathcal{S} is not minimal, and one can choose an x_0 such that

$$\deg P_z(\mathcal{S}) \leq \frac{\ell(y) - \ell(z) - 1}{2}, \quad \forall x_0 < z < y,$$

and

$$\deg P_{x_0}(\mathcal{S}) > \frac{\ell(y) - \ell(x_0) - 1}{2}.$$

(Cont'd next page)

7. Algorithm

(Algorithm cont'd)

By Proposition 6.4, there is an admissible set \mathcal{E}_1 which is dominated by \mathcal{S} and

$$h(\mathcal{E}_1) = h(\mathcal{S}) - gC'_{x_0}.$$

(3) One can identify the element g as follows. Let g_1 (resp. g_2) be the sum of terms in P_{x_0} of degree $> (\ell(y) - \ell(x_0) - 1)/2$ (resp. $> (\ell(y) - \ell(x_0))/2$), then

$$g = q^{-(\ell(y) - \ell(x_0))/2} g_1 + q^{(\ell(y) - \ell(x_0))/2} \bar{g}_2.$$

(4) Argue with \mathcal{E}_1 in place of \mathcal{S} and repeat steps (1) through (3).

It is clear that the algorithm terminates, and one can get a minimal admissible set and hence $P_{x,y}$ for all $x \leq y$.

8. Examples

- **Example 8.1.** Consider S_3 and $y = s_1 s_2 s_1$. Write $\ell(y, \sigma) = (\ell(y) - \ell(\pi(\sigma)) - 1)/2$. We have

\mathcal{S}	$\pi(\sigma)$	$d(\sigma)$	$\ell(y, \sigma)$
1, 1, 1	1	0	1
1, 1, s_1	s_1	0	1/2
1, s_2 , 1	s_2	0	1/2
s_1 , 1, 1	s_1	1	1/2
1, s_2 , s_1	$s_2 s_1$	0	0
s_1 , 1, s_1	1	1	1
s_1 , s_2 , 1	$s_1 s_2$	0	0
s_1 , s_2 , s_1	$s_1 s_2 s_1$		

So \mathcal{S} is not minimal. Take $x_0 = s_1$. Then $g = 1$ and $\mathcal{E}_1 = \mathcal{S} - \{(s_1, 1, 1), (s_1, 1, s_1)\} (= \mathcal{D}_0)$. Thus all $P_{x,y} = 1$ for $x \leq y$.

8. Examples

- **Example 8.2.** Consider S_4 and $y = s_2 s_1 s_3 s_2$. We have

\mathcal{S}	$\pi(\sigma)$	$d(\sigma)$	$\ell(y, \sigma)$
1, 1, 1, 1	1	0	3/2
1, 1, 1, s_2	s_2	0	1
1, 1, s_3 , 1	s_3	0	1
1, s_1 , 1, 1	s_1	0	1
s_2 , 1, 1, 1	s_2	1	1
1, 1, s_3 , s_2	$s_3 s_2$	0	1/2
1, s_1 , 1, s_2	$s_1 s_2$	0	1/2
s_2 , 1, 1, s_2	1	1	3/2
1, s_1 , s_3 , 1	$s_1 s_3$	0	1/2
s_2 , 1, s_3 , 1	$s_2 s_3$	0	1/2
s_2 , s_1 , 1, 1	$s_2 s_1$	0	1/2
1, s_1 , s_3 , s_2	$s_1 s_3 s_2$	0	0
s_2 , 1, s_3 , s_2	$s_2 s_3 s_2$	0	0
s_2 , s_1 , 1, s_2	$s_2 s_1 s_2$	0	0
s_2 , s_1 , s_3 , 1	$s_2 s_1 s_3$	0	0
s_2 , s_1 , s_3 , s_2			

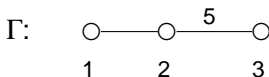
8. Examples

■ Example 8.2-cont'd

Hence \mathcal{S} is minimal and $P_{x,y} = P_x(\mathcal{S})$. We see that all $P_{x,y} = 1$ except

$$P_{1,y} = P_{s_2,y} = 1 + q.$$

■ Example 8.3. Consider H_3 . Its Coxeter graph is



Take $y = s_3 s_2 s_1 s_3 s_2 s_3$, then

$$\begin{aligned} P_{s_3 s_2 s_3, y} &= P_{s_3 s_2, y} = P_{s_2 s_3, y} = P_{s_2, y} = 1 + q, \\ P_{s_3, y} &= P_{1, y} = 1 + q + q^2, \end{aligned}$$

and all others are 1.