

# Tutorial: Economic Aspects of Social Networks (part II)

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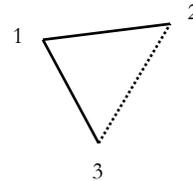
## Outline

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- Empirical regularities of social networks
- Capturing these features in models
- Strategic aspects of network formation
- Strategic interactions on networks

## Illustration of Concepts

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- **Degree:**
  - *With* link 23, degrees are (2, 2, 2)
  - *Without* link 23, degrees are (2, 1, 1)
- **Clustering:** What is the probability of link 23?
- **Diameter:**
  - *With* link 23, diameter = 1
  - *Without* link 23, diameter = 2

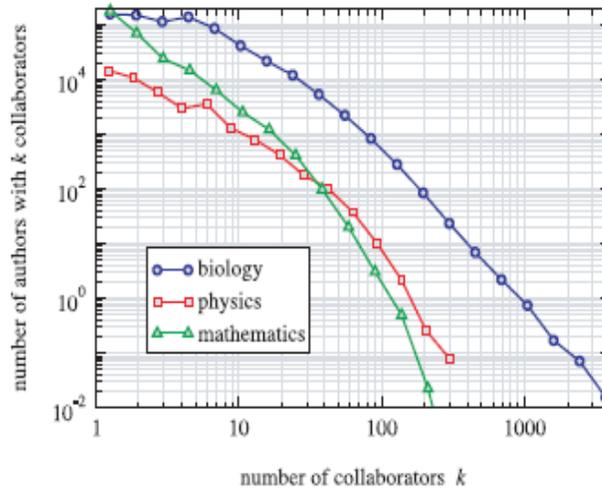
## Social Networks Share Many Features

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1. **Degree distributions:** Heavy tails, “scale-free”
  - Many more nodes with very high and very low degrees relative to what one would find in a completely random network

# Degree Distributions

Co-Authorship Data (Newman and Grossman)



# Features of Social Networks

1. **Degree distributions:** Heavy tails, “scale-free”
2. **High clustering** – “cliquishness”
  - Proportion of triads out of possible triads in a (sub)graph
  - Any two of a given node’s neighbors are likely to themselves be neighbors

Movie Actors	Math Co-authors	Physics Co-authors	WWW
0.79	0.43	0.15	0.11

## Features of Social Networks

1. **Degree distributions:** Heavy tails, “scale-free”
2. **High clustering** – “cliquishness”
3. **Low diameter** and average path length

	Movie Actors	Math Co-authors	Physics Co-authors	WWW
Average	3.7	7.6	5.9	3.1
Diameter		27	20	

## Features of Social Networks

1. **Degree distributions:** “Scale-free” – heavy tails
2. **High clustering** – “cliquishness”
3. **Low diameter** and average path length
4. **Assortativity**
  - Positive correlation in the degrees of linked nodes
  - Special to socially-generated networks

## Assortativity

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- Correlation in Degree (Newman 2003):
  - Socially generated networks:
    - .12 math co-authorship, .13 biology, .36 physics
    - .09 emails
    - .21 film actors
  - Technologically generated networks
    - -.19 internet
    - -.003 power grid
    - -.23 neural network

## Many other regularities

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- Negative relationship between degree and local clustering
- Homophily
- Strength of weak ties (Granovetter)
- Structural holes (Burt)
- Many more to be found. Which are the “right” ones?



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## Goals

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- Develop models with several goals in mind
  - Consistent with observed regularities
  - Based on reasonable micro-level assumptions
  - Analytically tractable
  - Yield novel insights

## A Network Formation Model (Jackson, Rogers)

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- Agents indexed by date of birth  $T = \{1, 2, 3, \dots\}$
- Upon birth, agent meets  $m_r$  others uniformly at random
- Search those neighborhoods to find  $m_n$  more nodes
  - Like entering at a random web page and following links
- Link to a given node if net utility is positive (probability  $p$ )

## Key Intuitions

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- Network-based meetings: likely to find nodes with many links
  - » variation on scale-free: growth partly proportional to size
- Random aspect: not entirely scale-free
  - » fit lower tail too
- Local search: may connect to two nodes that are already linked
  - » high clustering

## Key Intuitions (cont.)

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- Search aspect generates hub-like nodes
  - » low diameter
- Randomness connects different neighborhoods
  - » even lower diameter: order  $\ln(n)/\ln(\ln(n))$
- Nodes enter sequentially and connect to existing nodes
  - » Assortativity: age is correlated with both degree and links

## Mean-field approximation

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- Network formation is stochastic and path-dependent
  - Difficult to analyze directly
- Use a deterministic continuous-time system
  - All changes occur at the mean rate of underlying stochastic system
- Apply, e.g., to expected changes in degrees of nodes
- Can analytically solve for steady state

## Important Parameters

$$r = m_r/m_n$$

- Ratio of the number of links formed at random vs. by network-based meetings

$$m = p(m_r+m_n)$$

- Average number of out-links per node

## (In-)Degree Distribution

Expected increase in the in-degree of a node  $i$  at time  $t$  is roughly

$$p \left( \frac{m_r}{t} + [d_i \left( \frac{m_r}{t} \right)] \left[ \frac{m_n}{(m_r m)} \right] \right)$$

prob found at random (points to  $m_r/t$ )  
 prob found through search (points to the search term)  
 prob linked to given found (points to  $p$ )  
 number of neighbors (points to  $d_i$ )  
 prob a neighbor is entry point (points to  $m_r/t$ )  
 prob found from search of neighborhoods of entry points (points to  $m_n/(m_r m)$ )

## Theorem: Degree Distribution

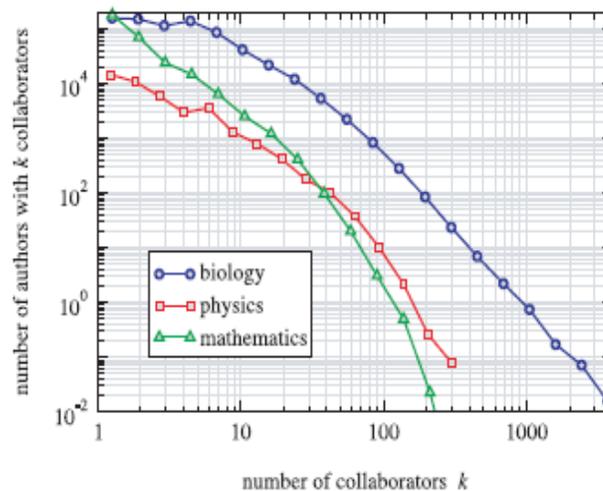
The in-degree distribution of the mean-field process has a degree distribution with cdf

$$F(d) = 1 - [r\mu/(d+r\mu)]^{1+r}$$

- Approximates a power distribution for large  $d$
- Lower tail is thinner

## Degree Distributions

Co-Authorship Data (Newman and Grossman)



## Sketch of Proof

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1. Mean-field:

$$\partial d_i(t)/\partial t = E[\text{change}(d_i(t))]; \quad d_i(0) = 0$$

2. Solve the differential equation:

$$d_i(t) = rm(t/i)^{1/(1+r)} - rm$$

3.  $1-F_t(d)$  = proportion of nodes with degree  $> d$  at time  $t$

- Let  $i^*(d)$  be such that  $d_{i^*(d)}(t) = d$
- $\Rightarrow 1-F_t(d) = i^*(d)/t$

4. Solve for  $i^*(d)$  from Step 2 and plug in

## Degree distribution shifts

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- Consider a distribution  $F$  with parameters  $(m, r)$  and a distribution  $F'$  with parameters  $(m', r')$ . Then:
  - If  $r' = r$  and  $m' > m$ , then  $F'$  strictly **FOSD**  $F$ 
    - I.e.,  $F'(d) < F(d)$  for all  $d$
  - If  $m' = m$  and  $r' < r$ , then  $F'$  is a strict **MPS** of  $F$

## Simple implications for efficiency

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- If utility to a node (information, employment, etc.) is increasing in degree, then average utility increases with  $m$ .
- If utility to a node is concave in degree, then average utility increases with  $r$ .

## Implications for Diffusion: SIS Model

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- Stochastic process where behavior/state of agents change depending on behavior/states of neighbors
  - Change from “S” to “I” with probability proportional to # of neighbors in state “I”
  - Revert from “I” to “S” at random
- Originally used to model the spreading of disease (common cold, computer virus)
- Can also model behavior
  - Adopting a new behavior (buying new product, coordination game)
  - Myopic best-response dynamics, may revert to status quo

## The model

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- Fix the network. Diffusion occurs in discrete time.
- $\theta_i$  is proportion of  $i$ 's infected neighbors
- "S"  $\rightarrow$  "I": probability  $v(d_i\theta_i)$
- "I"  $\rightarrow$  "S": probability  $0 < \delta < 1$
- Set  $\lambda = v/\delta$
  
- $\rho(d) =$  average infection rate of degree- $d$  nodes
- $\rho = \sum_d \rho(d)F(d) =$  average infection rate of society
- $\theta = \sum_d d\rho(d)F(d) / m$ 
  - Probability of linking to an infected node

## Results for SIS diffusion

(Jackson, Rogers)

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Proposition: There is positive steady-state infection iff  $\lambda > E[d]/E[d^2]$ .

Proposition: Suppose  $F$  is such that  $\theta^* > 0$ .

- If  $F'$  strictly FOSD  $F$ , then  $\theta^{*'} > \theta^*$  and  $\rho^{*'} > \rho^*$ .
- If  $F'$  is a strict MPS of  $F$ , then  $\theta^{*'} > \theta^*$ .
  - Note: overall infection rate  $\rho$  can go up or down

## Corollaries for Network Formation Model

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$$F(d) = 1 - [rm/(d + rm)]^{1+r}$$

- The minimal  $\lambda$  for positive steady-state infection is
  - 0 if  $r < 1$
  - $(r - 1)/2rm$  if  $r \geq 1$

Given  $F$  with  $(r, m)$  and  $F'$  with  $(r', m')$

- If  $r' = r$  and  $m' > m$ , then  $\theta^{*'} > \theta^*$  and  $\rho^{*'} > \rho^*$
- If  $m' = m$  and  $r' < r$ , then  $\theta^{*'} > \theta^*$

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## Main questions

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- What kind of structures are likely to form when linking choices are made by self-interested individuals?
- What are the welfare properties of these likely outcomes?
- What are the implications for processes occurring through the network?

## General setup

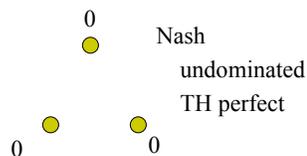
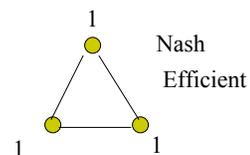
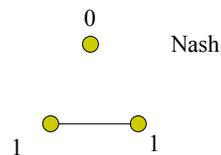
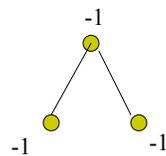
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- Agents  $N = \{1, 2, \dots, n\}$
- Network  $g$ , where  $g_{ij} = 1$  if  $i$  and  $j$  are connected, 0 otherwise
- Utility functions  $u_i(g) : G \rightarrow \mathbb{R}$
- Welfare criterion:  $U(g) = \sum_i u_i(g)$

## Solution concepts

- Consider the following game
  - Agents simultaneously announce potential links
  - Links that are mutually agreed are formed
  
- Nash equilibrium behaves poorly!
  - Multiplicity
    - E.g., the empty network is always an equilibrium
  - Cannot be easily refined away
    - Undominated strategies
    - Trembling-hand perfection

## Nash networks



## Pairwise stability

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- A network  $g$  is pairwise stable if
  - $g_{ij} = 1$  implies  $u_i(g) \geq u_i(g - ij)$ 
    - Nobody wants to delete a link (unilateral)
  - $g_{ij} = 0$  implies if  $u_i(g + ij) > u_i(g)$   
then  $u_j(g + ij) < u_j(g)$ 
    - No pair wants to add a link (bilateral)
- Game-form independent

## A basic “connections” model (Jackson, Wolinsky)

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- $u_i(g) = \sum_j \delta^{l_{ij}(g)} - cg_{ij}$ 
  - $c > 0$  link cost
  - $0 < \delta \leq 1$  decay parameter
  - $l_{ij}(g)$  length of shortest path between  $i$  and  $j$
- What will the set of pairwise stable networks look like?
- Will they be efficient?

## Efficient networks in the connections model

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Proposition: The efficient network is unique, and given by

- if  $c < \delta - \delta^2$  Complete
- if  $\delta - \delta^2 < c < \delta + (n-2)/2 * \delta^2$  Star
- if  $\delta + (n-2)/2 * \delta^2 < c$  Empty

## Pairwise stability in the connections model

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PS networks have the following properties

- There is at most one component
  - If  $c < \delta - \delta^2$  the only PS network is Complete
  - If  $\delta - \delta^2 < c < \delta$  then Star, and maybe others, are PS
  - If  $c > \delta$  then no agent has exactly one link
- Implication: PS networks are not necessarily efficient
- When  $c$  is slightly bigger than  $\delta$  then no PS can be efficient

## General tension between PS and Efficiency

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- PS is a mild requirement for stability
  - Not necessarily robust to many kinds of deviations
- What if we also consider transfers among players with the hope that we can induce efficient networks to be PS?
- This turns out to not be generally possible!

## General framework

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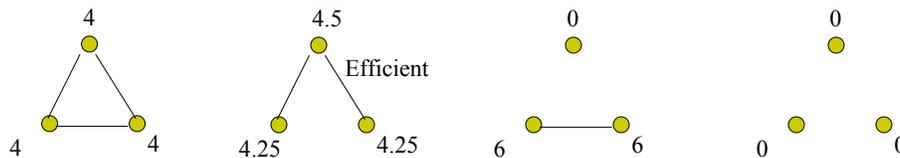
Transfers  $t:G \rightarrow \mathbb{R}^n$  such that  $\sum_i t_i(g) = 0$

- Require transfers to be
  - “Component Balanced” whenever utilities are component-based
    - Cannot make net transfers across components
  - Satisfy “Equal Treatment of Equals”
    - If two agents have the same set of neighbors, and always generate the same value to other agents, they must receive the same transfer
- Augment pairwise stability to respect transfers:
  - $g_{ij} = 1$  implies  $u_i(g) + t_i(g) \geq u_i(g - ij) + t_i(g - ij)$
  - $g_{ij} = 0$  implies if  $u_i(g + ij) + t_i(g + ij) > u_i(g) + t_i(g)$   
then  $u_j(g + ij) + t_j(g + ij) < u_j(g) + t_j(g)$

## Main result (Jackson, Wolinsky)

Proposition: There exist component-based utilities such that every network that is stable relative to a given transfer rule satisfying CB and ETE is not efficient.

## Sketch of proof



- Try to design  $t$  so that Eff network is PS
- ETE implies all  $t_i = 0$  for Complete & Empty
- ETE+CB imply  $t_i = 0$  for one-link network
- Take Eff network with agent 2 in center
  - ETE implies  $t_1 = t_3$
  - Need  $t_1 = t_3 \geq -.25$  (else add link 13)
  - Need  $t_2 \geq 1.5$  (else he severs a link)
  - But then  $t_1 + t_2 + t_3 > 0$  – a contradiction

## An “islands-connections” model (Jackson, Rogers)

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- Agents are grouped on “islands”
  - K islands
  - J agents per island
  
- Linking costs:
  - C for inter-island link, c for intra-island link
  - $C > c > 0$
  
- Benefit truncation
  - No benefit when distance is greater than D

## Small-world features

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- Proposition: Let  $c < \delta - \delta^2$  and  $C < \delta + (J - 1)\delta^2$ .  
Any network that is PS and/or Eff is such that:
- Islands are fully intra-connected
  - Diameter is no greater than  $D+1$
  - Clustering is greater than  $(J-1)(J-2)/(J^2K^2)$

## Many other aspects of strategic network formation

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- Strong stability
- Extensive form games
- Dynamic processes of link addition/deletion
- Directed networks, weighted networks
- A huge variety of specific utility models

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## Semi-anonymous games on graphs

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- Agents and graph given exogenously
- Agents choose binary action  $x_i = 0$  or  $1$
- Utility has the form  $u_d(x_i, m)$ 
  - Depends on degree  $d$ , own action, and
  - $m$  = number of neighbors choosing  $1$
- Goal: analyze equilibrium properties

## Two examples

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- Coordination game
  - $u_d(1, m) = a \cdot m / d$
  - $u_d(0, m) = b \cdot (d - m) / d$ 
    - Action 1 is optimal iff  $m/d \geq b/(a+b)$
- Best-shot game
  - $u_d(1, m) = 1 - c$  ( $0 < c < 1$ )
  - $u_d(0, m) = 1$  if  $m \geq 1$ ,  $0$  otherwise
    - Action 1 is optimal iff  $m = 0$

## General strategic features

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- A game has Strategic Complements if for all  $d$  and for all  $m \geq m'$ ,
  - $u_d(1,m) - u_d(0,m) \geq u_d(1,m') - u_d(0,m')$
  - Ex: coordination game
  
- Strategic Substitutes if “ $\leq$ ”
  - Ex: best-shot game

## Equilibrium properties under SC and SS

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- In both cases optimal behavior is characterized by “threshold” strategies
  - There is a function  $t(d)$  such that
    - SC:  $x_i = 1$  iff  $m \geq t(d)$
    - SS:  $x_i = 0$  iff  $m \geq t(d)$
  
- Games of SC: There always exists a pure strategy equilibrium, and the set of equilibria form a lattice
  
- Games of SS: There may not exist a pure strategy equilibrium

## Comparative statics w.r.t. graph

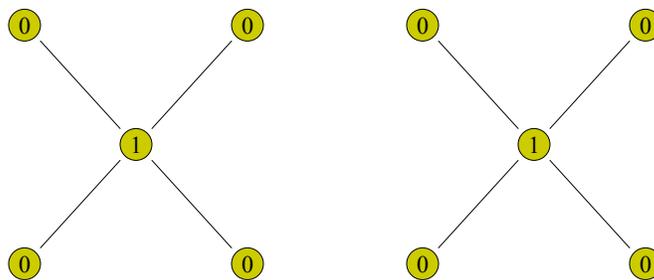
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Proposition: Take a game of SC with a non-increasing threshold  $t(d)$ .

- For every equilibrium  $x$  of  $g$  and for every  $g'$  larger than  $g$ , there exists an equilibrium  $x'$  of  $g'$  where  $x' \geq x$ .
  
- With SS, adding links can have hard-to-predict consequences
  - The “opposite” result of getting lower actions does not hold

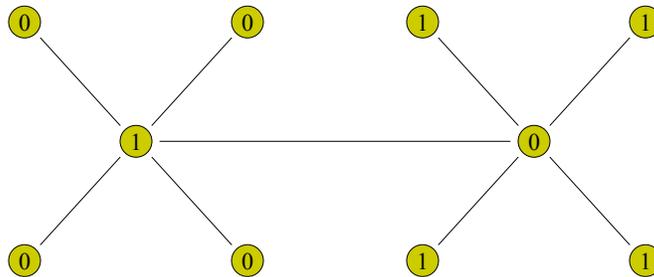
## An equilibrium of the best-shot game

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## Adding link causes many agents to switch

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## Result for best-shot game

(Galeotti et al)

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Proposition: Consider the best-shot game on  $g$  and an equilibrium  $x$  of  $g+ij$ .

- Either  $x$  is an equilibrium of  $g$ , or
- There exists an equilibrium of  $g$  in which a strict superset of agents chooses 1

□ “Adding links decreases actions”

## Other processes on networks

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- Semi-anonymous games in the configuration model
- Diffusion
- Information transmission (job openings)
- Opinion formation
- Learning (both Bayesian and non-Bayesian)
- Evolutionary dynamics
- Simultaneous evolution of links and actions

## Concluding remarks

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- Incorporating network elements is essential for many modeling applications
- The range of settings and analyses is very broad
- Many open questions
- Current research is becoming increasingly interdisciplinary
  - E.g., designing incentives in peer-to-peer applications