

# CHAPTER 6.

## FURTHER RESULTS

- 6.1. Propositions D-H.
- 6.2. Effectivity.
- 6.3. A Refutation.

### 6.1. Propositions D-H.

Our treatment of Propositions A,B,C culminated with Theorems 5.9.9, 5.9.11, and 5.9.12 at the end of Chapter 5.

In this section, we consider five Propositions D-H that have the same metamathematical properties as Propositions A,B,C. We will also consider some variants of Propositions D-H that do not share these properties, or whose status is left open.

Recall the main theorems of Chapter 5 (in section 5.9), which are Theorems 5.9.9, 5.9.11, and 5.9.12. Examination of the proofs of these three Theorems reveal that Theorem 5.9.11 with  $1\text{-Con}(\text{SMAH})$  is the key. If  $\text{ACA}'$  proves the equivalence of a statement with  $1\text{-Con}(\text{SMAH})$  then all of the other properties provided by these three Theorems quickly follow.

Accordingly, we establish these same three Theorems for Propositions D-H by showing that they are also each equivalent to  $1\text{-Con}(\text{SMAH})$  over  $\text{ACA}'$ .

We begin with Proposition D (see below), which is a sharpening of Proposition B. Proposition D immediately implies Propositions A-C over  $\text{RCA}_0$ .

Note that Propositions A-C are based on ELG. Examination of the proof of Proposition B in Chapter 4 shows that we can separately weaken the conditions on  $f, g$  in different ways. Also, we can place an inclusion condition on the starting set  $A_1$ . As usual, we use  $||$  for the sup norm, or max. This results in Proposition D below.

DEFINITION 6.1.1. We say that  $f$  is linearly bounded if and only if  $f \in \text{MF}$ , and there exists  $d$  such that for all  $x \in \text{dom}(f)$ ,

$$f(x) \leq d|x|.$$

We let LB be the set of all linearly bounded  $f$ .

DEFINITION 6.1.2. We say that  $g$  is expansive if and only if  $g \in MF$ , and there exists  $c > 1$  such that for all but finitely many  $x \in \text{dom}(f)$ ,

$$c|x| \leq g(x)$$

We let EXPN be the set of all expansive  $g$ .

Recall the definitions of MF, SD (Definition 1.1.2), and ELG, EVSD (Definitions 2.1, 2.2).

PROPOSITION D. Let  $f \in LB \cap EVSD$ ,  $g \in EXPN$ ,  $E \subseteq N$  be infinite, and  $n \geq 1$ . There exist infinite  $A_1 \subseteq \dots \subseteq A_n \subseteq N$  such that

- i) for all  $1 \leq i < n$ ,  $fA_i \subseteq A_{i+1} \cup gA_{i+1}$ ;
- ii)  $A_1 \cap fA_n = \emptyset$ ;
- iii)  $A_1 \subseteq E$ .

Note that  $ELG \subseteq LB \cap EVSD \cap EXPAN$ , and so Proposition D immediately implies Proposition B.

Proposition D is the strongest Proposition that we prove in this book (from large cardinals).

Recall that Propositions A-C are official statements of BRT. More accurately, Proposition B is really an infinite collection of statements of BRT.

Proposition D not a statement (or statements) of BRT for two reasons.

- a. There is no common set of functions used for  $f, g$  (asymmetry).
- b. The set  $E$  is used as data, rather than just  $f, g$ .

Features a,b both suggest very natural expansions of BRT. Feature a suggests "mixed BRT", where one uses several classes of functions instead of just one. One can go further and use several classes of sets as well.

Feature b in Proposition D suggests another very natural expansion of BRT. In BRT, we consider statements of the form

given functions there are sets such that  
 a given Boolean relation holds between the sets  
 and their images under the functions.

We can expand BRT with

given functions and sets there are sets such that  
 a given Boolean relation holds between the sets  
 and their images under the functions.

We will not pursue such expansions of BRT in this book.

We remark that feature  $b$  can be removed (in some contexts  
 such as here) by introducing a new function  $h$  and asserting  
 that  $A_1 \subseteq hN$  (obviously  $hN = \text{rng}(h)$ ).

We now prove Proposition D in SMAH<sup>+</sup> by adapting the proof of  
 Proposition B in SMAH<sup>+</sup> given in section 4.2.

We fix  $f, g, E$  as given by Proposition D. Analogously to  
 section 4.2, we let  $f$  be  $p$ -ary,  $g$  be  $q$ -ary. We fix an  
 integer  $b \geq 1$  such that for all  $x \in N^p$  and  $y \in N^q$ ,

i. if  $|x|, |y| > b$  then

$$\begin{aligned} |x| < f(x) &\leq b|x|. \\ (1 + 1/b)|y| &\leq g(y). \end{aligned}$$

ii. if  $|x| \leq b$  then  $f(x) \leq b^2$ .

Note how our inequalities are weaker than those used in  
 section 4.2.

We also fix  $n \geq 1$  and a strongly  $p^{n-1}$ -Mahlo cardinal  $\kappa$ .

The first place in section 4.2 that needs to be modified is  
 at Lemma 4.2.2. Here we must use the given infinite set  $E \subseteq N$ .

LEMMA 4.2.2'. There exist infinite sets  $E \supseteq E_0 \supseteq E_1 \supseteq \dots$   
 indexed by  $N$ , such that for all  $i \geq 0$ ,  $\varphi \in AF(L)$ ,  $\text{lth}(\varphi) \leq$   
 $i$ , and increasing partial  $h_1, h_2: V(L) \rightarrow N$  adequate for  $\varphi$  with  
 $\text{rng}(h_1), \text{rng}(h_2) \subseteq E_i$ , we have  $\text{Sat}(M, \varphi, h_1) \leftrightarrow \text{Sat}(M, \varphi, h_2)$ .

Proof: See the proof of Lemma 4.2.2. QED

Lemma 4.2.3 do not involve our inequalities  $i, ii$ , and therefore require no modification.

We need to sharpen Lemma 4.2.4 for later purposes, since we do not have an upper bound for  $g$ . We use the  $\#$  notation that was introduced much later just before Lemma 4.2.16.

LEMMA 4.2.4'. Let  $\varphi \in AS(L^*)$ .  $Sat(M^*, \varphi)$  if and only if  $\varphi \in T$ .  $<^*$  is a linear ordering on  $N^*$ . Let  $n \geq 0$ ,  $t \in CT(L^*)$ ,  $\#(t) \leq n$ . Then  $t < c_{n+1} \in T$ .

Proof: For the first claim, see the proof of Lemma 4.2.4. For the last claim, let  $i = lth(t < c_{n+1})$ . The unique increasing bijection  $h: V(L) \rightarrow E_i$  has  $Val(M, t', h) < h(v_{n+1})$ , where  $t'$  is the result of replacing each  $c_i$  by  $v_i$ , using the indiscernibility of  $E_i$ . Argue as before. QED

Lemmas 4.2.5 - 4.2.8 do not involve our inequalities  $i, ii$ , and therefore require no modification.

We sharpen Lemma 4.2.9 for later purposes, since we do not have an upper bound for  $g$ .

LEMMA 4.2.9'. These definitions of  $<^{**}$ ,  $+^{**}$ ,  $f^{**}$ ,  $g^{**}$  are well defined. Let  $t \in CT(L^{**})$ ,  $\#(t) \leq \alpha$ . Then  $t <^{**} c_{\alpha+1}^{**}$ .

Proof: Use Lemma 4.2.4' and the proof of Lemma 4.2.9. QED

Lemmas 4.2.10' - 4.2.14' do not involve our inequalities  $i, ii$ .

We need to weaken Lemma 4.2.15, in light of our inequalities  $i, ii$ .

LEMMA 4.2.15'. Let  $x_1, \dots, x_p, y_1, \dots, y_q \in N^{**}$ , where  $|x_1, \dots, x_p|, |y_1, \dots, y_q| >^{**} b^\wedge$ . Then

$$\begin{aligned} |x_1, \dots, x_p| <^{**} f^{**}(x_1, \dots, x_p) &\leq^{**} b|x_1, \dots, x_p|. \\ (1 + 1/b)|y_1, \dots, y_q| &\leq^{**} g^{**}(y_1, \dots, y_q). \end{aligned}$$

If  $|x_1, \dots, x_p| \leq^{**} b^\wedge$  then  $f(x_1, \dots, x_p) \leq^{**} b^{2^\wedge}$ .

Proof: See the proof of Lemma 4.2.15. QED

We aim for a modification of the crucial well foundedness given by Lemma 4.2.19. This was stated using all elements of  $N^{**}$ . In other words, for all terms in  $CT(L^{**})$ . We cannot

establish such a well foundedness result in the present setting for all terms in  $CT(L^{**})$ . We have weakened the inequalities for  $f^{**}, g^{**}$  too much.

However, we can establish this well foundedness result for the restricted class of terms,  $CT(L^{**} \setminus g)$  consisting of all closed terms of  $L^{**}$  in which  $g$  does not appear.

LEMMA 4.2.16'. Let  $t \in CT(L^{**})$ .  $\#(t) = -1 \leftrightarrow \text{Val}(M^{**}, t)$  is standard. Suppose  $\#(t) = c_\alpha$ . Then  $c_\alpha^{**} \leq \text{Val}(M^{**}, t) <^{**} c_{\alpha+1}^{**}$ . Let  $s \in CT(L^{**} \setminus g)$ . Suppose  $\#(s) = c_\alpha$ . There exists a positive integer  $d$  such that  $c_\alpha^{**} \leq^{**} \text{Val}(M^{**}, s) <^{**} dc_\alpha^{**} <^{**} c_{\alpha+1}^{**}$ .

Proof: For the equivalence in the first claim, see the proof of Lemma 4.2.16. For the remaining claims, use induction on  $s, t$ , Lemmas 4.2.4', 4.2.9', 4.2.15', and the proof of Lemma 4.2.16. QED

Lemmas 4.2.17, 4.2.18 do not involve our inequalities  $i, ii$ , and therefore require no modification.

DEFINITION 6.1.3. It is convenient to write  $VCT(L^{**} \setminus g)$  for the set of values of terms in  $CT(L^{**} \setminus g)$ .

DEFINITION 6.1.4. Let  $s$  be a rational number. We write  $<_s^{**}$  for the relation on  $VCT(L^{**} \setminus g)$  given by  $x <_s^{**} y \leftrightarrow sx <^{**} Y$ .

LEMMA 4.2.19'. Let  $s$  be a rational number  $> 1$ . There exists  $k \geq 1$  such that for all  $x_1 <_s^{**} x_2 <_s^{**} \dots <_s^{**} x_k$ , we have  $2x_1 <^{**} x_k$ .

Proof: See the proof of Lemma 4.2.19. QED

Lemma 4.2.20 has to be weakened as follows.

LEMMA 4.2.20'. Let  $s$  be a rational number  $> 1$ . The relation  $<_s^{**}$  on  $VCT(L^{**} \setminus g)$  is transitive, irreflexive, and well founded.

Proof: We adapt the proof of Lemma 4.2.20 with the following modification. In the fourth paragraph,  $d \in \mathbb{N} \setminus \{0\}$  is fixed such that  $\text{Val}(M^{**}, t) <^{**} dc_\alpha^{**}$ , using Lemma 4.2.16. Here we use Lemma 4.2.16' under the assumption that  $t \in VCT(L^{**} \setminus g)$ . QED

DEFINITION 6.1.5. Let  $s = 1 + 1/2b$  for using Lemma 4.2.20'.

LEMMA 4.2.21'. There is a unique set  $W$  such that  $W = \{x \in \text{VCT}(L^{**} \setminus g) \cap \text{nst}(M^{**}) : x \notin g^{**}W\}$ . For all  $\alpha < \kappa$ ,  $c_\alpha^{**} \notin \text{rng}(f^{**}), \text{rng}(g^{**})$ . In particular, each  $c_\alpha^{**} \in W$ .

Proof: Note that  $g^{**}: \text{NST}(M^{**})^q \rightarrow \text{NST}(M^{**})$ , but  $g^{**}: (\text{VCT}(L^{**} \setminus g) \cap \text{nst}(M^{**}))^q \rightarrow \text{VCT}(L^{**} \setminus g) \cap \text{nst}(M^{**})$  may be false. So we regard  $g^{**}$  as a partial function from  $(\text{VCTM}(L^{**} \setminus g) \cap \text{nst}(M^{**}))^q$  into  $\text{VCT}(L^{**} \setminus g) \cap \text{nst}(M^{**})$ . Note that  $g^{**}$  is strictly dominating from  $\text{nst}(M^{**})$  into  $\text{nst}(M^{**})$ , in the sense of  $<_s^{**}$ , by 4.2.15'. Since  $<_s^{**}$  is well founded on  $\text{VCT}(L^{**} \setminus g) \cap \text{nst}(M^{**})$ , we can apply the Complementation Theorem for Well Founded Relations, proved in section 1.3 to obtain the first claim.

For the second claim, write  $c_\alpha^{**} = f^{**}(x_1, \dots, x_p)$ . By Lemma 4.2.15', each  $x_i <^{**} c_\alpha^{**}$ . By Lemma 4.2.18,  $f^{**}(x_1, \dots, x_p) <^{**} c_\alpha^{**}$ . This is a contradiction. The same argument applies to  $g^{**}$ .

The third claim follows immediately from the second claim.  
QED

Lemma 4.2.22 - Theorem 4.2.26, Corollary 4.2.27, go through using the present  $W \subseteq \text{VCT}(L^{**} \setminus g) \cap \text{nst}(M^{**})$ , instead of the  $W \subseteq \text{nst}(M^{**})$  in section 4.2. We have shown the following.

THEOREM 6.1.1. Proposition D is provable in SMAH<sup>+</sup>. For fixed arity of  $f$  and fixed  $n \geq 1$ , Proposition D is provable in SMAH.

We now adapt section 4.4 to Proposition D. We redefine the  $p, q, b$ -structures,  $p, q, b; r$ -structures,  $p, q, b; n, r$ -special structures,  $p, q, b; r$ -types,  $p, q, b; n, r$ -special types, to take into account the weaker inequalities now placed on  $f, g$ . Specifically, clauses 4, 5 in the definition of  $p, q, b$ -structure should now read

4'.  $f^*$  obeys the above two inequalities for membership in  $\text{LB}(p, b) \cap \text{EVSD}(p, b)$  given above right after we introduced Proposition D, internally in  $M^*$ .

5'.  $g^*$  obeys the above two inequalities for membership in  $\text{EXP}(q, b)$ , given above right after we introduced Proposition D, internally in  $M^*$ .

These modified notions are written with '.

The entire development of section 4.4 goes through without modification until we arrive at Theorem 4.4.11.

THEOREM 4.4.11'. Proposition D is provable in  $ACA' + 1-Con(MAH)$ .

Proof: We argue in  $ACA' + 1-Con(MAH)$ . Let  $p, q, b, n \geq 1$ , and  $f \in LB(p, b) \cap EVSD(p, b)$ ,  $g \in EXPN(q, b)$ . Let  $r$  be given by Lemma 4.4.10'. By Ramsey's theorem for  $2r$ -tuples in  $ACA'$ , we can find a  $p, q, b; r$ -structure'  $M = (N, 0, 1, <, +, f, g, c_0, c_1, \dots)$ , where  $c_0, c_1, \dots \in E$ . Let  $\tau$  be its  $p, q, b; r$ -type'. By Lemma 4.4.10',  $\tau$  is a  $p, q, b, n, r$ -special' type. By Lemma 4.4.2,  $M$  is a  $p, q, b; r; n$ -special' structure. Let  $D_1 \subseteq \dots \subseteq D_n \subseteq N$ , where  $D_1 \subseteq \{c_0, c_1, \dots\} \subseteq E$ , and each  $fD_i \subseteq D_{i+1} \cup gD_{i+1}$ , and  $D_1 \cap fD_n = \emptyset$ . This is Proposition D, thus concluding the proof. QED

THEOREM 6.1.2.  $ACA'$  proves the equivalence of Proposition D and  $1-Con(MAH)$ ,  $1-Con(SMAH)$ .

Proof: This is immediate from Theorems 4.4.11', 5.9.11, and that Proposition D immediately implies Proposition B. QED

Recall that Proposition D is the strongest Proposition that we prove in this book (using large cardinals).

There are some natural variants of Proposition D, some of which are provable in  $RCA_0$ , and some of which are refutable.

PROPOSITION D[1]. Let  $f, g \in EVSD$ ,  $E \subseteq N$  be infinite, and  $n \geq 1$ . There exist infinite  $A_1 \subseteq \dots \subseteq A_n \subseteq N$  such that  
 i) for all  $1 \leq i < n$ ,  $fA_i \subseteq A_{i+1} \cup gA_{i+1}$ ;  
 ii)  $A_1 \cap fA_n = \emptyset$ ;  
 iii)  $A_1 \subseteq E$ .

Proposition D[1] is refutable in  $RCA_0$ . In fact, in section 6.3, we refute the following in  $RCA_0$ .

PROPOSITION  $\alpha$ . For all  $f, g \in SD \cap BAF$  there exist  $A, B, C \in INF$  such that

$$\begin{aligned} A \cup fA &\subseteq C \cup gB \\ A \cup fB &\subseteq C \cup gC. \end{aligned}$$

Note Proposition  $\alpha$  follows immediately from Proposition D[1], even without  $E$ . This is because from the former, we get

$$\begin{aligned}
& A \cup fA \subseteq B \cup gB \\
& A \cup fB \subseteq C \cup gC \\
& \quad B \subseteq C \\
& A \cup fA \subseteq C \cup gB.
\end{aligned}$$

Therefore Proposition D[1] is refutable in  $RCA_0$  even if we remove E.

However, we can use EVSD if we drop the inclusions on the A's.

PROPOSITION D[2]. Let  $f, g \in EVSD$ ,  $E \subseteq N$  be infinite, and  $n \geq 1$ . There exist infinite sets  $A_1, \dots, A_n \subseteq N$  such that

- i) for all  $1 \leq i < n$ ,  $fA_i \subseteq A_{i+1} \cup gA_{i+1}$ ;
- ii) for all  $1 \leq i \leq n$ ,  $A_1 \cap fA_n = \emptyset$ ;
- iii)  $A_1 \subseteq E$ .

The weakness in Proposition D[2] stems from the fact that we drop the tower condition, and use the same subscript twice on the right sides, and have no tower.

THEOREM 6.1.3. Proposition D[2] is provable in  $RCA_0$ .

Proof: Let  $f, g, E, n$  be as given. Let  $t \gg n \geq 1$ . By a straightforward combinatorial argument, for all  $t \geq 1$ , we can find an infinite  $E' \subseteq E$  such that

- a.  $f, g$  are strictly dominating on the elements of their respective domains whose sup norm is at least  $\min(E')$ .
- b. the values of all terms in  $f, g$  and elements of  $E'$ , using at most  $t$  applications of functions, and at least one application of a function, lie outside  $E'$ .

We now inductively define  $A_1, \dots, A_n$ . Set  $A_1 = E'$ . Suppose  $A_1, \dots, A_i$  have been defined for  $1 \leq i < n$ , where each  $A_j$  is an infinite subset of  $[\min(E'), \infty)$ . Set  $A_{i+1}$  to be the unique subset of  $fA_i$  such that  $fA_i \subseteq A_{i+1} \cup gA_{i+1}$ . This unique  $A_{i+1}$  exists by i) above and Lemma 3.3.3. Also  $A_{i+1}$  is infinite since  $fA_i$  is infinite (using a) above).

It is clear by the construction of the A's, that all elements of the  $fA_i$  and  $gA_i$  meet the criterion in b) above for  $t = n+1$ , so that their values lie outside  $E' = A_1$ . This establishes Proposition D[2] in  $RCA_0$ . QED



Continuing with our use of EVSD, it is natural to consider the following.

PROPOSITION D[3]. Let  $f, g \in \text{EVSD}$  and  $n \geq 1$ . There exist infinite sets  $A_1, \dots, A_n \subseteq \mathbb{N}$  such that

- i) for all  $1 \leq i < j, k \leq n$ ,  $fA_i \subseteq A_j \cup gA_k$ ;
- ii)  $A_1 \cap fA_n = \emptyset$ .

However, Proposition  $\alpha$  is an obvious consequence of Proposition D[3] even for the case  $n = 3$ . So Proposition D[3] is refutable in  $\text{RCA}_0$ .

PROPOSITION D[4]. Let  $f, g \in \text{EVSD}$ ,  $E \subseteq \mathbb{N}$  be infinite, and  $n \geq 1$ . There exist infinite sets  $A_1, \dots, A_n \subseteq \mathbb{N}$  such that

- i) for all  $1 \leq i < j, k \leq n$ ,  $fA_i \subseteq A_j \cup gA_k$ ;
- ii)  $A_1 \subseteq E$ .

THEOREM 6.1.4. Proposition D[4] is provable in  $\text{RCA}_0$ .

Proof: Let  $f, g, E$  be as given. Let  $m$  be such that  $f, g$  are strictly dominating on  $[m, \infty)$ . Let  $B$  be unique such that  $B \subseteq [m, \infty) \subseteq B \cup gB$ . Set  $A_1 = E \cap [m, \infty)$ ,  $A_2 = \dots = A_n = B$ . QED

PROPOSITION D[5]. Let  $f, g \in \text{EVSD}$  ( $\text{ELG}, \text{ELG} \cap \text{SD} \cap \text{BAF}$ ),  $E \subseteq \mathbb{N}$  be infinite, and  $n \geq 1$ . There exist  $A_1, \dots, A_n \subseteq \mathbb{N}$  such that

- i) for all  $1 \leq i < j, k \leq n$ ,  $fA_i \subseteq A_j \cup gA_k$ ;
- ii) for all  $1 \leq i \leq n$ ,  $A_i \cap E$  is infinite.

We do not know the status of Proposition D[5], other than it follows immediately from Proposition D.

We now present the remaining Propositions E, F that have the same metamathematical properties as Propositions A, B, C, D. These two propositions use  $\text{ELG} \cap \text{SD} \cap \text{BAF}$ .

DEFINITION 6.1.6. The powers of 2 are the integers  $1, 2, 4, 8, \dots$ . For  $E \subseteq \mathbb{N}$ , we write  $2^{(E)}$  for  $\{2^n : n \in E\}$ .

PROPOSITION E. For all  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  there exist  $A \subseteq B \subseteq C \subseteq \mathbb{N}$ , each containing infinitely many powers of 2, such that

$$\begin{aligned} fA &\subseteq B \cup gB \\ fB &\subseteq C \cup gC. \end{aligned}$$

PROPOSITION F. For all  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  there exist  $A \subseteq B \subseteq C \subseteq \mathbb{N}$ , each containing infinitely many powers of 2, such that

$$\begin{aligned} fA &\subseteq C \cup gB \\ fB &\subseteq C \cup gC. \end{aligned}$$

PROPOSITION G. For all  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  there exist  $A, B, C \subseteq \mathbb{N}$ , whose intersection contains infinitely many powers of 2, such that

$$\begin{aligned} fA &\subseteq C \cup gB \\ fB &\subseteq C \cup gC. \end{aligned}$$

PROPOSITION H. For all  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  there exist  $A, B, C \subseteq \mathbb{N}$ , where  $A \cap B$  contains infinitely many powers of 2, such that

$$\begin{aligned} fA &\subseteq C \cup gB \\ fB &\subseteq C \cup gC. \end{aligned}$$

Note that Propositions E-H are statements in BRT, where the BRT setting consists of "subsets of  $\mathbb{N}$  with infinitely many powers of 2", and  $\text{ELG} \cap \text{SD} \cap \text{BAF}$ . Propositions E, F, G immediately follow from Proposition D, using  $E = 2^{(\mathbb{N})}$ .

LEMMA 6.1.5. The following is provable in  $\text{RCA}_0$ .  $D \rightarrow E \rightarrow F \rightarrow G \rightarrow H$ .

Proof: For  $D \rightarrow E$ , let  $E = 2^{(\mathbb{N})}$ . For  $E \rightarrow F$ , use the derivation

$$\begin{aligned} fA &\subseteq B \cup gB \\ fB &\subseteq C \cup gC \\ B &\subseteq C \\ C \cap gB &= \emptyset \\ fA &\subseteq C \cup gB. \end{aligned}$$

$F \rightarrow G \rightarrow H$  is immediate. QED

We also consider two additional variants.

PROPOSITION E[1]. For all  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  there exist  $A, B, C \subseteq \mathbb{N}$ , whose intersection contains infinitely many powers of 2, such that

$$\begin{aligned} fA &\subseteq B \cup gB \\ fB &\subseteq C \cup gC. \end{aligned}$$

PROPOSITION G[1]. For all  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  there exist  $A, B, C \subseteq \mathbb{N}$ , each containing infinitely many powers of 2, such that

$$\begin{aligned} fA &\subseteq C \cup gB \\ fB &\subseteq C \cup gC. \end{aligned}$$

THEOREM 6.1.6. Proposition E[1] is provable in  $\text{RCA}_0$ .

Proof: Let  $f, g, E$  be as given. We follow the proof of Lemma 3.12.7. In the proof of Theorem 3.2.5, we can arrange that  $A \subseteq E$ . So in the proof of Lemma 3.12.7, we can assume that  $A \subseteq E$ . We also have  $A \subseteq B$ ,  $A \subseteq C$ . QED

We do not know the status of Proposition G[1], even if we use ELG instead of  $\text{ELG} \cap \text{SD} \cap \text{BAF}$ . Obviously, this follows from Proposition D with  $E = 2^{(\mathbb{N})}$ .

Until Theorem 6.1.10, we work in  $\text{RCA}_0$  and assume Proposition H.

LEMMA 6.1.7. For all  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  there exist infinite  $A, B, C \subseteq \mathbb{N}$  such that

$$\begin{aligned} fA &\subseteq C \cup gB \\ fB &\subseteq C \cup gC \\ A &\subseteq B, 2^{(\mathbb{N})}. \end{aligned}$$

Proof: Let  $f, g$  be as given. Let  $A, B, C$  be given by Proposition G. Replace  $A$  by  $A \cap B \cap 2^{(\mathbb{N})}$ , which is infinite. QED

LEMMA 6.1.8. The function  $f: \mathbb{N} \rightarrow \mathbb{N}$  given by  $f(n) = 1$  if  $n$  is a power of 2; 0 otherwise, lies in BAF.

Proof: Note that  $n$  is a power of 2 if and only if  $n = 2^{\log(n)}$ . QED

LEMMA 5.1.7'. Let  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$ . There exist  $f', g' \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  such that the following holds. Let  $S \subseteq \mathbb{N}$ .

- i)  $g'S = g(S^*) \cup 12S+2 \cup (f(S^*) \cap 2^{(\mathbb{N}+2)})$ .
- ii)  $f'S = f(S^*) \cup g'S \cup 12f(S^*)+2 \cup 2S^*+1 \cup 3S^*+1$ .

Proof: Let  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$ , where  $f: \mathbb{N}^p \rightarrow \mathbb{N}$  and  $g: \mathbb{N}^q \rightarrow \mathbb{N}$ . We define  $g': \mathbb{N}^{q+p} \rightarrow \mathbb{N}$  as follows. Let  $x_1, \dots, x_q, y_1, \dots, y_p \in \mathbb{N}$ .

case 1.  $x_1, \dots, x_q > y_1, \dots, y_p$ . Set  $g'(x_1, \dots, x_q, y_1, \dots, y_p) = g(x_1, \dots, x_q)$ .

case 2.  $y_1, \dots, y_p > x_1, \dots, x_q$  and  $f(y_1, \dots, y_p) \in 2^{(N+2)}$ . Set  $g'(x_1, \dots, x_q, y_1, \dots, y_p) = f(y_1, \dots, y_p)$ .

case 3. Otherwise. Set  $g'(x_1, \dots, x_q, y_1, \dots, y_p) = 12|x_1, \dots, x_q, y_1, \dots, y_p| + 2$ .

We define  $f': \mathbb{N}^{5p+q+p} \rightarrow \mathbb{N}$  as follows. Let  $x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p \in \mathbb{N}$ .

case a.  $|y_1, \dots, y_q, z_1, \dots, z_p| = |x_1, \dots, x_p| = |x_{p+1}, \dots, x_{2p}| = |x_{2p+1}, \dots, x_{3p}| = |x_{3p+1}, \dots, x_{4p}| = |x_{4p+1}, \dots, x_{5p}|$ . Set  $f'(x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p) = g'(y_1, \dots, y_q, z_1, \dots, z_p)$ .

case b.  $|y_1, \dots, y_q, z_1, \dots, z_p| = |x_1, \dots, x_p| = |x_{p+1}, \dots, x_{2p}| = |x_{2p+1}, \dots, x_{3p}| = |x_{3p+1}, \dots, x_{4p}| < \min(x_{4p+1}, \dots, x_{5p})$ . Set  $f'(x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p) = f(x_{4p+1}, \dots, x_{5p})$ .

case c.  $|y_1, \dots, y_{q+1}, z_1, \dots, z_p| = |x_1, \dots, x_p| = |x_{p+1}, \dots, x_{2p}| = |x_{2p+1}, \dots, x_{3p}| = |x_{4p+1}, \dots, x_{5p}| < \min(x_{3p+1}, \dots, x_{4p})$ . Set  $f'(x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p) = 12f(x_{3p+1}, \dots, x_{4p}) + 2$ .

case d.  $|y_1, \dots, y_q, z_1, \dots, z_p| = |x_1, \dots, x_p| = |x_{p+1}, \dots, x_{2p}| = |x_{3p+1}, \dots, x_{4p}| = |x_{4p+1}, \dots, x_{5p}| < \min(x_{2p+1}, \dots, x_{3p})$ . Set  $f'(x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p) = 2|x_{2p+1}, \dots, x_{3p}| + 1$ .

case e.  $|y_1, \dots, y_q, z_1, \dots, z_p| = |x_1, \dots, x_p| = |x_{2p+1}, \dots, x_{3p}| = |x_{3p+1}, \dots, x_{4p}| = |x_{4p+1}, \dots, x_{5p}| < \min(x_{2p+1}, \dots, x_{3p})$ . Set  $f'(x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p) = 3|x_{p+1}, \dots, x_{2p}| + 1$ .

case f. Otherwise. Set  $f'(x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p) = 2|x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p| + 1$ .

Note that in case 1,  $|x_1, \dots, x_q, y_1, \dots, y_p| = |x_1, \dots, x_q|$ , and in case 2,  $|x_1, \dots, x_q, y_1, \dots, y_p| = |y_1, \dots, y_p|$ . Also note that in cases a)-e),

$$\begin{aligned} |x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p| &= |y_1, \dots, y_q, z_1, \dots, z_p| \\ |x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p| &= |x_{4p+1}, \dots, x_{5p}| \\ |x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p| &= |x_{3p+1}, \dots, x_{4p}| \\ |x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p| &= |x_{2p+1}, \dots, x_{3p}| \\ |x_1, \dots, x_{5p}, y_1, \dots, y_q, z_1, \dots, z_p| &= |x_{p+1}, \dots, x_{2p}| \end{aligned}$$

respectively. Hence  $f', g' \in \text{ELG} \cap \text{SD} \cap \text{BAF}$ .

Let  $S \subseteq \mathbb{N}$ . From  $S$ , case 1 produces exactly  $g(S^*)$ . Case 2 produces exactly  $f(S^*) \cap 2^{(N+2)}$ . Case 3 produces exactly  $12S+2$ . This establishes i).

Case a) produces exactly  $g'S$ . Case b) produces exactly  $f(S^*)$ . Case c) produces exactly  $12f(S^*)+2$ . Case d) produces exactly  $2S^*+1$ . Case e produces exactly  $3S^*+1$ .

Case f) produces exactly  $2S^*+1$  since  $2\min(S)+1$  is not produced. This is because  $2\min(S)+1$  is produced from case f) if and only if all of the arguments are  $\min(S)$ , which can only happen under case a). This establishes ii). QED

LEMMA 6.1.9.  $12E+2$ ,  $6E$ ,  $2E+1 \cup 3E+1$ ,  $2^{(N+2)}$  are pairwise disjoint, with the sole exception of  $2E+1 \cup 3E+1$  and  $2^{(N+2)}$ .

Proof: Obviously,  $12E+2$ ,  $6E$ ,  $2E+1 \cup 3E+1$  are pairwise disjoint by divisibility considerations. Also  $12n+2 = 2m \rightarrow 6n+1 = 2^{m-1}$ , which is impossible for  $m \geq 3$ . QED

LEMMA 5.1.8'. Let  $f, g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$  and  $\text{rng}(g) \subseteq 6\mathbb{N}$ . There exist infinite  $A \subseteq B \subseteq C \subseteq \mathbb{N} \setminus \{0\}$  such that

- i)  $fA \cap 6\mathbb{N} \subseteq B \cup gB$ ;
- ii)  $fB \cap 6\mathbb{N} \subseteq C \cup gC$ ;
- iii)  $fA \cap 2\mathbb{N}+1 \subseteq B$ ;
- iv)  $fA \cap 3\mathbb{N}+1 \setminus 2^{(N+2)} \subseteq B$ ;
- v)  $fB \cap 2\mathbb{N}+1 \subseteq C$ ;
- vi)  $fB \cap 3\mathbb{N}+1 \setminus 2^{(N+2)} \subseteq C$ ;
- vii)  $C \cap gC = \emptyset$ ;
- viii)  $A \cap fB = \emptyset$ .

Proof: Let  $f, g$  be as given. Let  $f', g'$  be given by Lemma 5.1.7'. Let  $A, B, C \subseteq \mathbb{N}$  be given by Lemma 6.1.7 for  $f', g'$ . Then  $A, B, C$  are infinite, and

$$\begin{aligned} f'A &\subseteq C \cup g'B \\ f'B &\subseteq C \cup g'C \\ A &\subseteq B, 2^{(N)}. \end{aligned}$$

Since we can shrink  $A$  to any infinite subset, we will assume that  $A \subseteq 2^{(N+2)}$ .

Let  $n \in B$ . Then  $12n+2 \in g'B \cap f'B$ , and so  $12n+2 \in C \cup g'C$ . Now  $12n+2 \notin C$  by  $C \cap g'B = \emptyset$ . Hence  $12n+2 \in g'C$ . Therefore  $12n+2 \in 12C+2$ . Hence  $n \in C$ . So we have established that  $A \subseteq B \subseteq C$ .

We now verify all of the required conditions i)-viii) above using the three sets  $A^*, B^*, C^*$ .

Firstly note that  $A^* \subseteq B^* \subseteq C^* \subseteq N \setminus \{0\}$ . To see this, first observe that  $\min(A) \geq \min(B) \geq \min(C)$ . Now let  $n \in A^*$ . Then  $n \in B \wedge n > \min(A) \geq \min(B)$ . Hence  $n \in B^*$ . Thus  $A^* \subseteq B^*$ . The same argument establishes  $B^* \subseteq C^*$ .

We now claim that  $A^* \cap f(B^*) = \emptyset$ . Let  $n \in A^*$ ,  $n \in f(B^*)$ . Then  $n \in f(B^*) \cap 2^{(N+2)}$ ,  $n \in g'B$ ,  $n \in C$ . This is a contradiction.

Next we claim that  $C^* \cap g(C^*) = \emptyset$ . This follows from  $C \subseteq C^*$ ,  $g(C^*) \subseteq g'C$ , and  $C \cap g'C = \emptyset$ .

Now we claim that  $f(A^*) \cap 6N \subseteq B^* \cup g(B^*)$ . To see this, let  $n \in f(A^*) \cap 6N$ . Then  $n \in f'A$ ,  $n \in C \cup g'B$ .

case 1.  $n \in C$ . Now  $12n+2 \in g'C$  and  $12n+2 \in 12f(A^*)+2 \subseteq f'A$ . Since  $C \cap g'C = \emptyset$ , we have  $12n+2 \notin C$ . Also  $12n+2 \in C \cup g'B$ . Hence  $12n+2 \in g'B$ . Therefore  $12n+2 \in 12B+2$ , and so  $n \in B$ . Since  $n \in f(A^*)$  and  $f$  is strictly dominating, we have  $n > \min(A) \geq \min(B)$ . Hence  $n \in B^*$ .

case 2.  $n \in g'B$ . Since  $n \in 6N$ ,  $n \in g(B^*)$ . This establishes the claim.

Next we claim that  $f(B^*) \cap 6N \subseteq C^* \cup g(C^*)$ . To see this, let  $n \in f(B^*) \cap 6N$ . Then  $n \in f'B$ . Hence  $n \in C \cup g'C$ .

case 1'.  $n \in C$ . Since  $n \in f(B^*)$  and  $f$  is strictly dominating, we have  $n > \min(B) \geq \min(C)$ . Hence  $n \in C^*$ .

case 2'.  $n \in g'C$ . Since  $n \in 6N$ , we have  $n \in g(C^*)$ . This establishes the claim.

Now we claim that  $f(A^*) \cap 2N+1$ ,  $f(A^*) \cap 3N+1 \setminus 2^{(N+2)} \subseteq B^*$ . To see this, let  $n \in f(A^*)$ ,  $n \in 2N+1 \cup 3N+1$ ,  $n \notin 2^{(N+2)}$ . Note that  $n \notin \text{rng}(g')$ . Also,  $n \in f'A$ ,  $n \in C \cup g'B$ . Hence  $n \in C$ ,  $12n+2 \in g'C$ ,  $12n+2 \notin C$ . Now  $12n+2 \in 12f(A^*)+2 \subseteq f'A \subseteq C \cup g'B$ ,  $12n+2 \in g'B$ ,  $n \in B$ . Since  $f$  is strictly dominating,  $n > \min(A) \geq \min(B)$ , and so  $n \in B^*$ .

Finally we claim that  $f(B^*) \cap 2N+1$ ,  $f(B^*) \cap 3N+1 \setminus 2^{(N+2)} \subseteq C^*$ . To see this, let  $n \in f(B^*)$ ,  $n \in 2N+1 \cup 3N+1$ ,  $n \notin 2^{(N+2)}$ . Note that  $n \notin \text{rng}(g')$ . Also,  $n \in f'B$ ,  $n \in C \cup g'C$ . Hence  $n \in C$ ,  $12n+2 \in g'C$ ,  $12n+2 \notin C$ . Now  $12n+2 \in 12f(B^*)+2 \subseteq f'B \subseteq$

$C \cup g'C$ . Hence  $12n+2 \in g'C$ ,  $n \in C$ . Since  $f$  is strictly dominating,  $n > \min(B) \geq \min(C)$ , and so  $n \in C^*$ . QED

The proof of 1-Con(SMAH) from Proposition C given in Chapter 5 is strictly modular, in that we can start with Lemma 5.1.8 instead of Proposition C.

Here we repeat the proof in Chapter 5 using Lemma 5.1.8' instead of Lemma 5.1.8. However, Lemma 5.1.8' is slightly weaker than Lemma 5.1.8, because of the weakened clauses iv) and vi), where we use  $3N+1 \setminus 2^{(N+2)}$  instead of  $3N+1$ .

So we need to identify the few places at which we use  $3N+1$  and make sure that we can get away with  $3N+1 \setminus 2^{(N+2)}$  instead.

By examination of the proofs, we obtain the following series of slightly weakened Lemmas from the end of sections 5.1 - 5.5. Finally, we show that we obtain Lemma 5.6.20 without modification.

LEMMA 5.2.12'. Let  $r \geq 3$  and  $g \in \text{ELG} \cap \text{SD} \cap \text{BAF}$ , where  $\text{rng}(g) \subseteq 48N$ . There exists  $(D_1, \dots, D_r)$  such that

- i)  $D_1 \subseteq \dots \subseteq D_r \subseteq N \setminus \{0\}$ ;
- ii)  $|D_1| = r$  and  $D_r$  is finite;
- iii) for all  $x < y$  from  $D_1$ ,  $x \uparrow < y$ ;
- iv) for all  $1 \leq i \leq r-1$ ,  $48\alpha(r, D_i; 1, r) \subseteq D_{i+1} \cup gD_{i+1}$ ;
- v) for all  $1 \leq i \leq r-1$ ,  $2\alpha(r, D_i; 1, r) + 1$ ,  $3\alpha(r, D_i; 1, r) + 1 \setminus 2^{(N+2)} \subseteq D_{i+1}$ ;
- vi)  $D_r \cap gD_r = \emptyset$ ;
- vii)  $D_1 \cap \alpha(r, D_2; 2, r) = \emptyset$ ;
- viii) Let  $1 \leq i \leq \beta(2r)$ ,  $x_1, \dots, x_{2r} \in D_1$ ,  $y_1, \dots, y_r \in \alpha(r, D_2)$ , where  $(x_1, \dots, x_r)$  and  $(x_{r+1}, \dots, x_{2r})$  have the same order type and  $\min$ , and  $y_1, \dots, y_r \leq \min(x_1, \dots, x_r)$ . Then  $t[i, 2r](x_1, \dots, x_r, y_1, \dots, y_r) \in D_3 \leftrightarrow t[i, 2r](x_{r+1}, \dots, x_{2r}, y_1, \dots, y_r) \in D_3$ .

LEMMA 5.3.18'. There exists a countable structure  $M = (A, <, 0, 1, +, -, \cdot, \uparrow, \log, E, c_1, c_2, \dots)$  such that the following holds.

- i)  $(A, <, 0, 1, +, -, \cdot, \uparrow, \log)$  satisfies  $\text{TR}(\Pi_1^0, L)$ ;
- ii)  $E \subseteq A \setminus \{0\}$ ;
- iii) The  $c_n$ ,  $n \geq 1$ , form a strictly increasing sequence of nonstandard elements in  $E \setminus \alpha(E; 2, < \infty)$  with no upper bound in  $A$ ;
- iv) Let  $r, n \geq 1$ ,  $t(v_1, \dots, v_r)$  be a term of  $L$ , and  $x_1, \dots, x_r \leq c_n$ . Then  $t(x_1, \dots, x_r) < c_{n+1}$ ;
- v)  $2\alpha(E; 1, < \infty) + 1$ ,  $3\alpha(E; 1, < \infty) + 1 \setminus 2^{(A+2)} \subseteq E$ ;

- vi) Let  $r \geq 1$ ,  $a, b \in \mathbb{N}$ , and  $\varphi(v_1, \dots, v_r)$  be a quantifier free formula of  $L$ . There exist  $d, e, f, g \in \mathbb{N} \setminus \{0\}$  such that for all  $x_1 \in \alpha(E; 1, < \infty)$ ,  $(\exists x_2, \dots, x_r \in E) (x_2, \dots, x_r \leq ax_1 + b \wedge \varphi(x_1, \dots, x_r)) \leftrightarrow dx_1 + e \notin E \leftrightarrow fx_1 + g \in E$ ;
- vii) Let  $r \geq 1$ ,  $p \geq 2$ , and  $\varphi(v_1, \dots, v_{2r})$  be a quantifier free formula of  $L$ . There exist  $a, b, d, e \in \mathbb{N} \setminus \{0\}$  such that the following holds. Let  $n \geq 1$  and  $x_1, \dots, x_r \in \alpha(E; 1, < \infty) \cap [0, c_n]$ . Then  $(\exists y_1, \dots, y_r \in E) (y_1, \dots, y_r \leq \uparrow p(|x_1, \dots, x_r|) \wedge \varphi(x_1, \dots, x_r, y_1, \dots, y_r)) \leftrightarrow a \text{CODE}(c_{n+1}; x_1, \dots, x_r) + b \notin E \leftrightarrow d \text{CODE}(c_{n+1}; x_1, \dots, x_r) + e \in E$ . Here CODE is as defined just before Lemma 5.3.11;
- viii) Let  $k, n, m \geq 1$ , and  $x_1, \dots, x_k \leq c_n < c_m$ , where  $x_1, \dots, x_k \in \alpha(E; 1, < \infty)$ . Then  $\text{CODE}(c_m; x_1, \dots, x_k) \in E$ ;
- ix) Let  $r \geq 1$  and  $t(v_1, \dots, v_{2r})$  be a term of  $L$ . Let  $i_1, \dots, i_{2r} \geq 1$  and  $y_1, \dots, y_r \in E$ , where  $(i_1, \dots, i_r)$  and  $(i_{r+1}, \dots, i_{2r})$  have the same order type and  $\min$ , and  $y_1, \dots, y_r \leq \min(c_{i_1}, \dots, c_{i_r})$ . Then  $t(c_{i_1}, \dots, c_{i_r}, y_1, \dots, y_r) \in E \leftrightarrow t(c_{i_{r+1}}, \dots, c_{i_{2r}}, y_1, \dots, y_r) \in E$ .

Lemma 5.4.12 uses  $2\alpha(E; 1, < \infty) + 1$ ,  $3\alpha(E; 1, < \infty) + 1 \subseteq E$ . However, we only have  $3\alpha(E; 1, < \infty) + 1 \setminus 2^{(A+2)} \subseteq E$ . So it suffices to augment the displayed derivation in Lemma 5.4.12 with the second derivation

$$\begin{aligned}
& t(x_1, \dots, x_k) < c_{n+1}. \\
& 2c_{n+1} + t(x_1, \dots, x_k) + 3, 3c_{n+1} + t(x_1, \dots, x_k) + 2 \in \alpha(E; 1, < \infty). \\
& 3(2c_{n+1} + t(x_1, \dots, x_k) + 2) + 1, 2(3c_{n+1} + t(x_1, \dots, x_k) + 3) + 1 \in E. \\
& 6c_{n+1} + 3t(x_1, \dots, x_k) + 7, 6c_{n+1} + 2t(x_1, \dots, x_k) + 7 \in E. \\
& (6c_{n+1} + 3t(x_1, \dots, x_k) + 7) - (6c_{n+1} + 2t(x_1, \dots, x_k) + 7) = \\
& t(x_1, \dots, x_k) \in E - E.
\end{aligned}$$

provided we verify that

$$3(2c_{n+1} + t(x_1, \dots, x_k)) + 1 \notin 2^{(A+2)} \vee 3(2c_{n+1} + t(x_1, \dots, x_k) + 2) + 1 \notin 2^{(A+2)}.$$

This is evident, since any two powers of 2 that are  $\geq 4$  cannot differ by 6.

LEMMA 5.4.17'. There exists a countable structure  $M = (A, <, 0, 1, +, -, \cdot, \uparrow, \log, E, c_1, c_2, \dots)$ , and terms  $t_1, t_2, \dots$  of  $L$ , where for all  $i$ ,  $t_i$  has variables among  $v_1, \dots, v_{i+8}$ , such that the following holds.

- i)  $(A, <, 0, 1, +, -, \cdot, \uparrow, \log)$  satisfies  $\text{TR}(\Pi_1^0, L)$ ;



- ii)  $E \subseteq A \setminus \{0\}$ ;
- iii) The  $c_n$ ,  $n \geq 1$ , form a strictly increasing sequence of nonstandard elements in  $E \setminus \alpha(E; 2, < \infty)$  with no upper bound in  $A$ ;
- iv) Let  $r, n \geq 1$  and  $t(v_1, \dots, v_r)$  be a term of  $L$ , and  $x_1, \dots, x_r \leq c_n$ . Then  $t(x_1, \dots, x_r) < c_{n+1}$ ;
- v)  $2\alpha(E; 1, < \infty) + 1, 3\alpha(E; 1, < \infty) + 1 \setminus 2^{(A+2)} \subseteq E$ ;
- vi) Let  $k, n \geq 1$  and  $R$  be a  $c_n$ -definable  $k$ -ary relation. There exists  $y_1, \dots, y_8 \in E \cap [0, c_{n+1}]$  such that  $R = \{(x_1, \dots, x_k) \in E^k \cap [0, c_n]^k : t_k(x_1, \dots, x_k, y_1, \dots, y_8) \in E\}$ ;
- vii) Let  $r \geq 1$  and  $\varphi(v_1, \dots, v_{2r})$  be a formula of  $L(E)$ . Let  $1 \leq i_1, \dots, i_{2r} < n$ , where  $(i_1, \dots, i_r)$  and  $(i_{r+1}, \dots, i_{2r})$  have the same order type and the same min. Let  $y_1, \dots, y_r \in E$ ,  $y_1, \dots, y_r \leq \min(c_{i_1}, \dots, c_{i_r})$ . Then  $\varphi(c_{i_1}, \dots, c_{i_r}, y_1, \dots, y_r)^{c_{-n}} \Leftrightarrow \varphi(c_{i_{r+1}}, \dots, c_{i_{2r}}, y_1, \dots, y_r)^{c_{-n}}$ .

LEMMA 5.5.8'. There exists a countable structure  $M^* = (A, <, 0, 1, +, -, \cdot, \uparrow, \log, E, c_1, c_2, \dots, X_1, X_2, \dots)$ , where for all  $i \geq 1$ ,  $X_i$  is the set of all  $i$ -ary relations on  $A$  that are  $c_n$ -definable for some  $n \geq 1$ ; and terms  $t_1, t_2, \dots$  of  $L$ , where for all  $i$ ,  $t_i$  has variables among  $x_1, \dots, x_{i+8}$ , such that the following holds.

- i)  $(A, <, 0, 1, +, -, \cdot, \uparrow, \log)$  satisfies  $TR(\Pi^0_1, L)$ ;
- ii)  $E \subseteq A \setminus \{0\}$ ;
- iii) The  $c_n$ ,  $n \geq 1$ , form a strictly increasing sequence of nonstandard elements of  $E \setminus \alpha(E; 2, < \infty)$  with no upper bound in  $A$ ;
- iv) For all  $r, n \geq 1$ ,  $\uparrow r(c_n) < c_{n+1}$ ;
- v)  $2\alpha(E; 1, < \infty) + 1, 3\alpha(E; 1, < \infty) + 1 \setminus 2^{(A+2)} \subseteq E$ ;
- vi) Let  $k, n \geq 1$  and  $R$  be a  $c_n$ -definable  $k$ -ary relation. There exists  $y_1, \dots, y_8 \in E \cap [0, c_{n+1}]$  such that  $R = \{(x_1, \dots, x_k) \in E^k \cap [0, c_n]^k : t_k(x_1, \dots, x_k, y_1, \dots, y_8) \in E\}$ ;
- vii) Let  $k \geq 1$ ,  $m \geq 0$ , and  $\varphi$  be an  $E$  formula of  $L^*(E)$  in which  $R$  is not free, where all first order variables free in  $\varphi$  are among  $x_1, \dots, x_{k+m+1}$ . Then  $x_{k+1}, \dots, x_{k+m+1} \in E \rightarrow (\exists R)(\forall x_1, \dots, x_k \in E)(R(x_1, \dots, x_k) \Leftrightarrow (x_1, \dots, x_k \leq x_{k+m+1} \wedge \varphi))$ ;
- viii) Let  $r \geq 1$ , and  $\varphi(x_1, \dots, x_{2r})$  be an  $E$  formula of  $L^*(E)$  with no free second order variables. Let  $1 \leq i_1, \dots, i_{2r}$ , where  $(i_1, \dots, i_r)$  and  $(i_{r+1}, \dots, i_{2r})$  have the same order type and the same min. Let  $x_1, \dots, x_r \in E$ ,  $x_1, \dots, x_r \leq \min(c_{i_1}, \dots, c_{i_r})$ . Then  $\varphi(c_{i_1}, \dots, c_{i_r}, x_1, \dots, x_r) \Leftrightarrow \varphi(c_{i_{r+1}}, \dots, c_{i_{2r}}, x_1, \dots, x_r)$ .

Lemma 5.6.2 involves reproving a weak form of Lemma 5.4.12 using a related construction. Here  $3\alpha(E; 1, < \infty) + 1 \subseteq E$  can also be replaced by  $3\alpha(E; 1, < \infty) + 1 \setminus 2^{(A+2)}$ , also by the same method.

In the remainder of section 5.6, we do not use  $3\alpha(E;1,<\infty)+1\setminus 2^{(A+2)} \subseteq E$ . Hence we obtain Lemma 5.6.20. We have proved the following.

THEOREM 6.1.10. ACA' proves that each of Propositions A-H are equivalent to Con(SMAH).

Proof: We have completed the proof that ACA' proves Proposition H implies 1-Con(SMAH). The result follows by Lemmas 5.9.11 and 6.1.5. QED